## Course Review for Midterm II

Cpt S 223
Fall 2009

## Midterm Exam 2

- When: Friday (11/20) 10:10-11am
- Where: in class
- Closed book, closed notes
- Comprehensive
- With primary focus on the following topics: priority queues, hashing and disjoint sets.
- Other topics covered before Midterm exam 1 could be included although indirectly.
- Material for preparation:
- Lecture slides \& in-class notes
- Homeworks \& program assignments
- Weiss book


## Priority Queues

- Binary heap, Binomial heaps
- Definitions:
- Binary heap
- Structure property:
- complete binary tree except the last level (filled breadth-first left to right)
- Heap order property
- Min heap : each node's value is less than or equal to its children
- Binomial heap
- Structure property:
- A forest of binomial trees (similar to binary representation)
- Heap order property:
- Min heap: within each binomial tree, same heap order like in binary heap


## Implementation

- Binary heap
- Tree structure can be implemented as an array
- Where nodes are stored in breadth-first order
- Children of node at $\mathrm{A}[i]$ are at: $\mathrm{A}[2 i]$ and $A[2 i+1]$
- Conversion procedure: Array to tree, tree to array
- Binomial heap
- Array of pointers to each binomial tree
- $\log \mathrm{n}$ binomial tree pointers
- Know relationships between binomial heap, binomial trees and their properties


## Run-times for each heap operation

Two main techniques: percolateUp and percolateDown

|  | Insert | DeleteMin | Merge |
| :--- | :--- | :--- | :--- |
| Binary heap | O(lg n) worst-case, <br> O(1) using <br> buildheap | O(lg n) | O(n) |
| Binomial Heap | O(lg n) worst-case <br> O(1) for insertion in <br> sequence | O(lg n) | O(lg n) |

Other operations:

- deleteMax()
- decreaseKey(p,v), increaseKey(p,v)
- remove(p)


## Union-Find data structure

- Union-find
- Supports two operations on disjoint sets:
- Union(a,b)
- Find(a)
- One application is equivalence class computation:
- Maximal subsets defined by equivalence relation
- Disjoint subsets
- Array implementation


## Steps in the Union (x, y)

Steps in Union $(x, y)$ :

1. EqClass $_{x}=$ Find $(x)$
2. EqClass $_{y}=$ Find $(y)$
3. EqClass $_{\mathrm{xy}}=$ UnionSets $\left(\right.$ EqClass $_{\mathrm{x}}$, EqClass $\left._{\mathrm{y}}\right)$

Equivalence class algorithm:

- Initially, put each element in a set of its own
- FOR EACH element pair (a,b):
- Check [a Rb=true]
- IFaRbTHEN
- Union(a,b)


## Variations in Union

- Simple Union (aka Arbitrary Union)
- Perform arbitrary linking of roots
- Smart Unions

1. Union-by-Rank

- Connect shorter tree under taller tree

2. Union-by-Size

- Connect smaller size tree under larger size tree (size(tree) = \#nodes in the tree)


## Variations of Find

- Simple Find
- Simply returns the root id without modifying the tree
- Smart Find
- Uses path compression
- Link all nodes along the path from $x$ to the root directly to the root
- Returns the root id, but the resulting tree could have been modified as a result of path compression


## Heuristics \& their Gains

|  | Worst-case run-time for $m$ operations |
| :---: | :---: |
| Arbitrary Union, Simple Find | $\mathrm{O}(\mathrm{m} \mathrm{n})$ |
| Union-by-size, Simple Find | $\mathrm{O}(\mathrm{m} \log \mathrm{n})$ |
| Union-by-rank, Simple Find | $\mathrm{O}(\mathrm{m} \log \mathrm{n})$ |
| Arbitrary Union, Path compression Find | $\begin{array}{\|rl\|} \hline \mathrm{O}(\mathrm{~m} \log \mathrm{n}) & \\ & \text { Extremely slow } \\ & \text { Growing functio } \\ \hline \end{array}$ |
| Union-by-rank, Path compression Find | O(m <br> Inv.Ackermann(m,n)) <br> $=O\left(m \log ^{*} n\right)$ |

## Hashing

- Hash functions
- purpose: string to integer, integer to integer
- Choice of a "good" hash functions
- Reduce chance of collision
- Relatively smaller key value
- Does not need huge hash table size
- Hash Tables use hash functions
- Hash table size should be a prime
- Load factor
- Measure to tell how crowded a hash table is
- Know algorithms \& analysis for the following
- Collision resolution by chaining
- Collision resolution by open-addressing
- Linear probing, quadratic probing
- Double hashing
- Rehashing
- Think of applications where hash tables could work out better than other data structures.


## General Tips

- Work out examples of basic operations
- Show steps whenever possible
- Questions mostly objective, but be prepared to see one or two subjective ones too
- Write less and to the point
- Don't leave questions totally unanswered
- If unable to solve, write at least your approach idea(s)

