

Lemma: Binomial Heap insertion sequence of  $n$  elements takes  $\mathcal{O}(n)$  time

Proof: Idea  $\Rightarrow$  For every new insertion, # comparisons = # position of least significant zero.  $(LSZ) + 1$

Bit Position:	4	3	2	1	0	Position of least significant zero (LSZ)
	0	0	0	0	0 <sup>↙</sup>	0
	0	0	0	0 <sup>↙</sup>	1	1
	0	0	0	1	0 <sup>↙</sup>	0
	0	0	0 <sup>↙</sup>	1	1	2
	0	0	1	0 <sup>↙</sup>	0 <sup>↙</sup>	0
	0	0	1	0 <sup>↙</sup>	1	1
	0	0	1	1	0 <sup>↙</sup>	0
	0	0 <sup>↙</sup>	1	1	1	3
	0	1	0	0	0 <sup>↙</sup>	0
	0	1	0	0 <sup>↙</sup>	1	1
	0	1	0 <sup>↙</sup>	1	1	2
	0	1	1	0	0 <sup>↙</sup>	0
	0	1	1	0 <sup>↙</sup>	1	1
	0	1	1	1	0 <sup>↙</sup>	0
	0 <sup>↙</sup>	1	1	1	1	4
Total <del>cost</del> $\approx$ 15.						
Cost $\propto 2(15)$						

LSZ cases:  
 don't decrease  $\rightarrow$  takes time  
 $\dots 0 \rightarrow 1$   
 $\dots 01 \rightarrow 2$   
 $\dots 011 \rightarrow 3$   
 $\dots 0111 \rightarrow 4$   
 $\vdots$   
 $0111 \dots 1 \rightarrow k$  time  
 $\underbrace{\hspace{2cm}}_{k-1}$

$\Rightarrow$  Total (amortized) cost of  $n$  successive inserts

$$= \frac{n}{2} \times 1 + \frac{n}{4} \times 2 + \frac{n}{8} \times 3 + \dots + \frac{n}{2^k} \times k$$

$$= \sum_{i=1}^k \frac{n}{2^i} \times i, \text{ where } k = \lg n$$

$$< \sum_{i=0}^{\infty} \frac{n}{2^i} \times i$$

$$= n \sum_{i=1}^{\infty} \frac{i}{2^i} = 2n$$

$\Rightarrow$  Total amortized cost of  $n$  successive inserts  $\leq 2n = \mathcal{O}(n)$

$\Rightarrow$  Amortized cost of each insert operation  $= \mathcal{O}(1)$