FIRST-ORDER LOGIC

CHAPTER 8
Outline

◊ Why FOL?
◊ Syntax and semantics of FOL
◊ Fun with sentences
◊ Wumpus world in FOL
Pros and cons of propositional logic

Propositional logic is **declarative**: pieces of syntax correspond to facts.

Propositional logic allows partial/disjunctive/negated information (unlike most data structures and databases).

Propositional logic is **compositional**: meaning of $B_{1,1} \land P_{1,2}$ is derived from meaning of $B_{1,1}$ and of $P_{1,2}$.

Meaning in propositional logic is **context-independent** (unlike natural language, where meaning depends on context).

Propositional logic has very limited expressive power (unlike natural language).

E.g., cannot say “pits cause breezes in adjacent squares” except by writing one sentence for each square.
First-order logic

Whereas propositional logic assumes world contains facts, first-order logic (like natural language) assumes the world contains

- **Objects**: people, houses, numbers, theories, Ronald McDonald, colors, baseball games, wars, centuries . . .
- **Relations**: red, round, bogus, prime, multistoried . . .
  brother of, bigger than, inside, part of, has color, occurred after, owns, comes between, . . .
- **Functions**: father of, best friend, third inning of, one more than, end of . . .
<table>
<thead>
<tr>
<th>Language</th>
<th>Ontological Commitment</th>
<th>Epistemological Commitment</th>
</tr>
</thead>
<tbody>
<tr>
<td>Propositional logic</td>
<td>facts</td>
<td>true/false/unknown</td>
</tr>
<tr>
<td>First-order logic</td>
<td>facts, objects, relations</td>
<td>true/false/unknown</td>
</tr>
<tr>
<td>Temporal logic</td>
<td>facts, objects, relations, times</td>
<td>true/false/unknown</td>
</tr>
<tr>
<td>Probability theory</td>
<td>facts</td>
<td>degree of belief</td>
</tr>
<tr>
<td>Fuzzy logic</td>
<td>facts + degree of truth</td>
<td>known interval value</td>
</tr>
</tbody>
</table>

Logics in general
Syntax of FOL: Basic elements

Constants  $KingJohn, 2, UCB,...$
Predicates  $Brother, >,...$
Functions  $Sqrt, LeftLegOf,...$
Variables  $x, y, a, b,...$
Connectives  $\land, \lor, \neg, \Rightarrow, \Leftrightarrow$
Equality  $=$
Quantifiers  $\forall, \exists$
Atomic sentences

Atomic sentence = \( \text{predicate}(\text{term}_1, \ldots, \text{term}_n) \)

or \( \text{term}_1 = \text{term}_2 \)

Term = \( \text{function}(\text{term}_1, \ldots, \text{term}_n) \)

or constant or variable

E.g., \( \text{Brother}(\text{KingJohn}, \text{RichardTheLionheart}) > (\text{Length}(\text{LeftLegOf}(\text{Richard})), \text{Length}(\text{LeftLegOf}(\text{KingJohn}))) \)
Complex sentences

Complex sentences are made from atomic sentences using connectives

\[ \neg S, \ S_1 \land S_2, \ S_1 \lor S_2, \ S_1 \Rightarrow S_2, \ S_1 \Leftrightarrow S_2 \]

E.g. \[ \text{Sibling}(\text{King John}, \text{Richard}) \Rightarrow \text{Sibling}(\text{Richard}, \text{King John}) \]
\[ >(1,2) \lor \leq (1,2) \]
\[ >(1,2) \land \neg >(1,2) \]
Sentences are true with respect to a model and an interpretation.

Model contains $\geq 1$ objects (domain elements) and relations among them.

Interpretation specifies referents for:
- Constant symbols $\rightarrow$ objects
- Predicate symbols $\rightarrow$ relations
- Function symbols $\rightarrow$ functional relations

An atomic sentence $\text{predicate}(\text{term}_1, \ldots, \text{term}_n)$ is true iff the objects referred to by $\text{term}_1, \ldots, \text{term}_n$ are in the relation referred to by $\text{predicate}$.
Models for FOL: Example

- Person R
  - Brother of J
  - Left leg
- Person J
  - Brother of R
  - Left leg
  - Crown on head
  - Person king

Chapter 8
Consider the interpretation in which

\textit{Richard} \rightarrow \text{Richard the Lionheart}

\textit{John} \rightarrow \text{the evil King John}

\textit{Brother} \rightarrow \text{the brotherhood relation}

Under this interpretation, \textit{Brother}(\textit{Richard, John}) is true just in case Richard the Lionheart and the evil King John are in the brotherhood relation in the model.
Models for FOL: Lots!

Entailment in propositional logic can be computed by enumerating models

We can enumerate the FOL models for a given KB vocabulary:

For each number of domain elements $n$ from 1 to $\infty$
  For each $k$-ary predicate $P_k$ in the vocabulary
    For each possible $k$-ary relation on $n$ objects
      For each constant symbol $C$ in the vocabulary
        For each choice of referent for $C$ from $n$ objects . . .

Computing entailment by enumerating FOL models is not easy!
Universal quantification

$$\forall \langle \text{variables} \rangle \langle \text{sentence} \rangle$$

Everyone at Berkeley is smart:
$$\forall x \ At(x, Berkeley) \Rightarrow Smart(x)$$

$$\forall x \ P$$ is true in a model $$m$$ iff $$P$$ is true with $$x$$ being each possible object in the model

Roughly speaking, equivalent to the conjunction of instantiations of $$P$$

$$(At(KingJohn, Berkeley) \Rightarrow Smart(KingJohn))$$
$$\land (At(Richard, Berkeley) \Rightarrow Smart(Richard))$$
$$\land (At(Berkeley, Berkeley) \Rightarrow Smart(Berkeley))$$
$$\land \ldots$$
A common mistake to avoid

Typically, $\Rightarrow$ is the main connective with $\forall$

Common mistake: using $\land$ as the main connective with $\forall$:

$$\forall x \ At(x, Berkeley) \land Smart(x)$$

means “Everyone is at Berkeley and everyone is smart”
Existential quantification

\[ \exists \langle \text{variables} \rangle \langle \text{sentence} \rangle \]

Someone at Stanford is smart:
\[ \exists x \ At(x, \text{Stanford}) \land \text{Smart}(x) \]

\[ \exists x \ P \] is true in a model \( m \) iff \( P \) is true with \( x \) being some possible object in the model

Roughly speaking, equivalent to the disjunction of instantiations of \( P \)

\[ \begin{align*}
(At(KingJohn, Stanford) \land \text{Smart}(KingJohn)) \\
\lor (At(Richard, Stanford) \land \text{Smart}(Richard)) \\
\lor (At(Stanford, Stanford) \land \text{Smart}(Stanford)) \\
\lor \ldots
\end{align*} \]
Another common mistake to avoid

Typically, $\land$ is the main connective with $\exists$

Common mistake: using $\Rightarrow$ as the main connective with $\exists$:

$$\exists x \ At(x, Stanford) \Rightarrow Smart(x)$$

is true if there is anyone who is not at Stanford!
Properties of quantifiers

\( \forall x \ \forall y \) is the same as \( \forall y \ \forall x \) \( \text{(why??)} \)

\( \exists x \ \exists y \) is the same as \( \exists y \ \exists x \) \( \text{(why??)} \)

\( \exists x \ \forall y \) is \textbf{not} the same as \( \forall y \ \exists x \)

\( \exists x \ \forall y \ Loves(x, y) \)
“There is a person who loves everyone in the world”

\( \forall y \ \exists x \ Loves(x, y) \)
“Everyone in the world is loved by at least one person”

Quantifier duality: each can be expressed using the other

\( \forall x \ \text{Likes}(x, \text{IceCream}) \quad \neg \exists x \ \neg \text{Likes}(x, \text{IceCream}) \)

\( \exists x \ \text{Likes}(x, \text{Broccoli}) \quad \neg \forall x \ \neg \text{Likes}(x, \text{Broccoli}) \)
Fun with sentences

Brothers are siblings
Fun with sentences

Brothers are siblings

\( \forall x, y \ \text{Brother}(x, y) \Rightarrow \text{Sibling}(x, y). \)

“Sibling” is symmetric
Fun with sentences

Brothers are siblings

\[ \forall x, y \ Brother(x, y) \Rightarrow Sibling(x, y). \]

“Sibling” is symmetric

\[ \forall x, y \ Sibling(x, y) \Leftrightarrow Sibling(y, x). \]

One’s mother is one’s female parent
Brothers are siblings

\[ \forall x, y \; \text{Brother}(x, y) \implies \text{Sibling}(x, y). \]

“Sibling” is symmetric

\[ \forall x, y \; \text{Sibling}(x, y) \iff \text{Sibling}(y, x). \]

One’s mother is one’s female parent

\[ \forall x, y \; \text{Mother}(x, y) \iff (\text{Female}(x) \land \text{Parent}(x, y)). \]

A first cousin is a child of a parent’s sibling
Fun with sentences

Brothers are siblings

\[ \forall x, y \; \text{Brother}(x, y) \Rightarrow \text{Sibling}(x, y). \]

“Sibling” is symmetric

\[ \forall x, y \; \text{Sibling}(x, y) \Leftrightarrow \text{Sibling}(y, x). \]

One’s mother is one’s female parent

\[ \forall x, y \; \text{Mother}(x, y) \Leftrightarrow (\text{Female}(x) \land \text{Parent}(x, y)). \]

A first cousin is a child of a parent’s sibling

\[ \forall x, y \; \text{FirstCousin}(x, y) \Leftrightarrow \exists p, ps \; \text{Parent}(p, x) \land \text{Sibling}(ps, p) \land \text{Parent}(ps, y) \]
Equality

\(term_1 = term_2\) is true under a given interpretation if and only if \(term_1\) and \(term_2\) refer to the same object.

E.g., \(1 = 2\) and \(\forall x \ (Sqrt(x), Sqrt(x)) = x\) are satisfiable.

\(2 = 2\) is valid.

E.g., definition of (full) \(Sibling\) in terms of \(Parent\):
\[
\forall x, y \ Sibling(x, y) \iff [\neg(x = y) \land \exists m, f \ \neg(m = f) \land \ Parent(m, x) \land Parent(f, x) \land Parent(m, y) \land Parent(f, y)]
\]
Interacting with FOL KBs

Suppose a wumpus-world agent is using an FOL KB and perceives a smell and a breeze (but no glitter) at \( t = 5 \):

\[
\text{Tell}(\text{KB}, \text{Percept}([\text{Smell}, \text{Breeze}, \text{None}], 5))
\]
\[
\text{Ask}(\text{KB}, \exists a \ \text{Action}(a, 5))
\]

I.e., does \( \text{KB} \) entail any particular actions at \( t = 5 \)?

Answer: \textit{Yes,} \{\( a/\text{Shoot} \}\} \leftarrow \text{substitution} \ (\text{binding list})

Given a sentence \( S \) and a substitution \( \sigma \),
\( S\sigma \) denotes the result of plugging \( \sigma \) into \( S \); e.g.,
\[
S = \text{Smarter}(x, y)
\]
\[\sigma = \{x/\text{Hillary}, y/\text{Bill}\}\]
\[
S\sigma = \text{Smarter}(\text{Hillary}, \text{Bill})
\]

\( \text{Ask}(\text{KB}, S) \) returns some/all \( \sigma \) such that \( \text{KB} \models S\sigma \)
Knowledge base for the wumpus world

“Perception”
\[
\forall b, g, t \ Percept([\text{Smell}, b, g], t) \Rightarrow \text{Smelt}(t)
\]
\[
\forall s, b, t \ Percept([s, b, \text{Glitter}], t) \Rightarrow \text{AtGold}(t)
\]

Reflex: \( \forall t \ \text{AtGold}(t) \Rightarrow \text{Action}(\text{Grab}, t) \)

Reflex with internal state: do we have the gold already?
\[
\forall t \ \text{AtGold}(t) \land \neg\text{Holding}(\text{Gold}, t) \Rightarrow \text{Action}(\text{Grab}, t)
\]

\(\text{Holding}(\text{Gold}, t)\) cannot be observed
\[ \Rightarrow \text{keeping track of change is essential} \]
Deducing hidden properties

Properties of locations:
\[ \forall x, t \; \text{At}(Agent, x, t) \land \text{Smelt}(t) \Rightarrow \text{Smelly}(x) \]
\[ \forall x, t \; \text{At}(Agent, x, t) \land \text{Breeze}(t) \Rightarrow \text{Breezy}(x) \]

Squares are breezy near a pit:

**Diagnostic** rule—infer cause from effect
\[ \forall y \; \text{Breezy}(y) \Rightarrow \exists x \; \text{Pit}(x) \land \text{Adjacent}(x, y) \]

**Causal** rule—infer effect from cause
\[ \forall x, y \; \text{Pit}(x) \land \text{Adjacent}(x, y) \Rightarrow \text{Breezy}(y) \]

Neither of these is complete—e.g., the causal rule doesn’t say whether squares far away from pits can be breezy

**Definition** for the \textit{Breezy} predicate:
\[ \forall y \; \text{Breezy}(y) \Leftrightarrow [\exists x \; \text{Pit}(x) \land \text{Adjacent}(x, y)] \]
Keeping track of change

Facts hold in situations, rather than eternally
E.g., Holding(Gold, Now) rather than just Holding(Gold)

Situation calculus is one way to represent change in FOL:
   Adds a situation argument to each non-eternal predicate
   E.g., Now in Holding(Gold, Now) denotes a situation

Situations are connected by the Result function
Result(a, s) is the situation that results from doing a in s
Describing actions I

“Effect” axiom—describe changes due to action
\( \forall s \ AtGold(s) \Rightarrow Holding(Gold, Result(Grab, s)) \)

“Frame” axiom—describe non-changes due to action
\( \forall s \ HaveArrow(s) \Rightarrow HaveArrow(Result(Grab, s)) \)

Frame problem: find an elegant way to handle non-change
   (a) representation—avoid frame axioms
   (b) inference—avoid repeated “copy-overs” to keep track of state

Qualification problem: true descriptions of real actions require endless caveats—what if gold is slippery or nailed down or …

Ramification problem: real actions have many secondary consequences—what about the dust on the gold, wear and tear on gloves, …
Describing actions II

Successor-state axioms solve the representational frame problem

Each axiom is “about” a predicate (not an action per se):

\[ P \text{ true afterwards} \iff \begin{cases} \text{an action made } P \text{ true} \\ \lor \text{ } P \text{ true already and no action made } P \text{ false} \end{cases}\]

For holding the gold:

\[ \forall a, s \quad Holding(Gold, Result(a, s)) \iff \begin{cases} (a = \text{Grab} \land AtGold(s)) \\ \lor (Holding(Gold, s) \land a \neq \text{Release}) \end{cases} \]
Making plans

Initial condition in KB:

\[ \text{At(Agent, [1, 1], S_0)} \]
\[ \text{At(Gold, [1, 2], S_0)} \]

Query: \( \text{Ask(KB, } \exists s \ \text{Holding(Gold, s)}) \)

i.e., in what situation will I be holding the gold?

Answer: \( \{s/\text{Result(Grab, Result(Forward, S_0))}\} \)

i.e., go forward and then grab the gold

This assumes that the agent is interested in plans starting at \( S_0 \) and that \( S_0 \) is the only situation described in the KB.
Making plans: A better way

Represent plans as action sequences \([a_1, a_2, \ldots, a_n]\)

\(\text{PlanResult}(p, s)\) is the result of executing \(p\) in \(s\)

Then the query \(\text{Ask}(\text{KB}, \exists p \ \text{Holding}(\text{Gold, PlanResult}(p, S_0)))\)
has the solution \(\{p/[\text{Forward, Grab}]\}\)

Definition of \(\text{PlanResult}\) in terms of \(\text{Result}\):
\[
\forall s \ \text{PlanResult}([], s) = s \\
\forall a, p, s \ \text{PlanResult}([a|p], s) = \text{PlanResult}(p, \text{Result}(a, s))
\]

Planning systems are special-purpose reasoners designed to do this type of inference more efficiently than a general-purpose reasoner
Summary

First-order logic:
- objects and relations are semantic primitives
- syntax: constants, functions, predicates, equality, quantifiers

Increased expressive power: sufficient to define wumpus world

Situation calculus:
- conventions for describing actions and change in FOL
- can formulate planning as inference on a situation calculus KB